Service-oriented Distributed Applications in the Future Internet: The Case for Interaction Paradigm Interoperability

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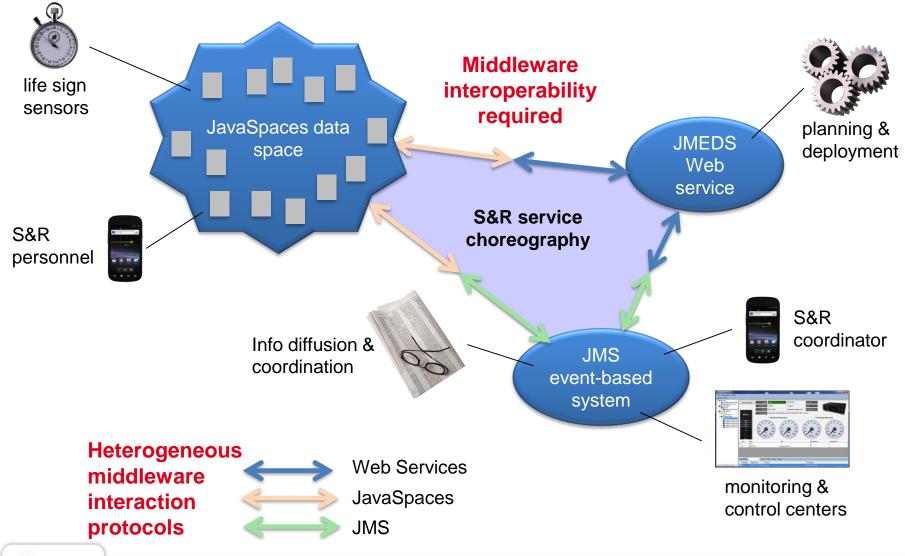
IFIP WG on Services-oriented Systems, Amsterdam, 11/10/2013







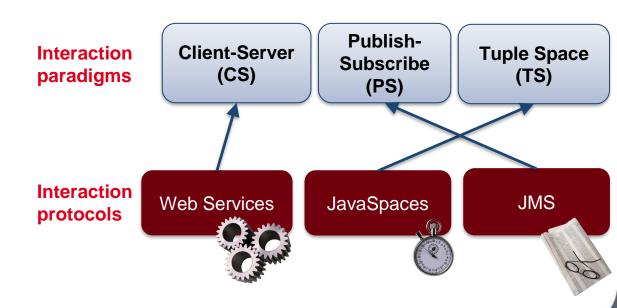
Service-oriented applications in the FI: Search & Rescue (S&R)



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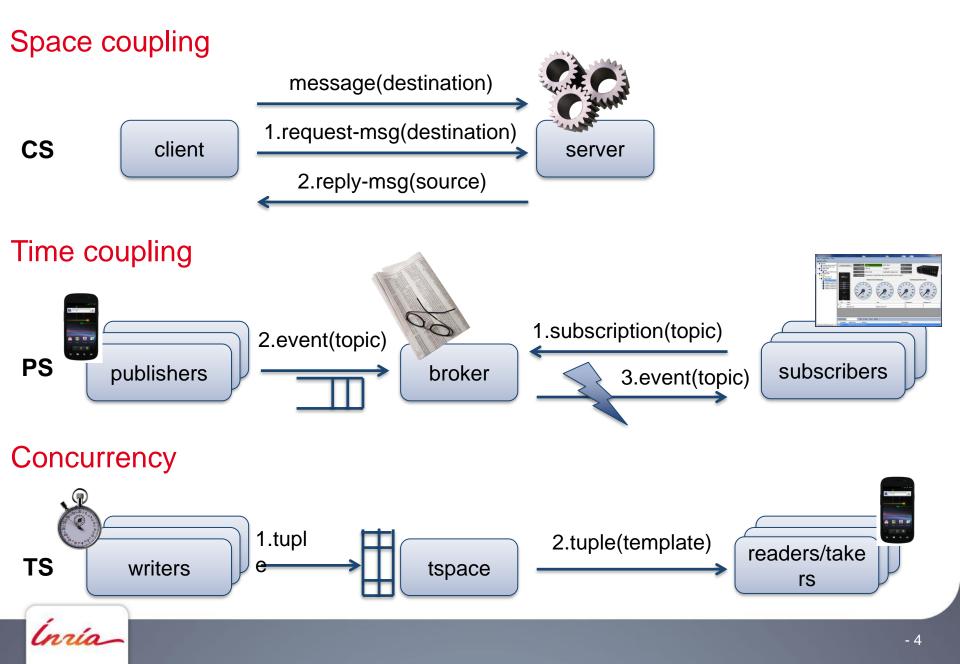
Interaction paradigms in middleware

- Look for comprehensive systematic solution to middleware interoperability
- Classify middleware implementations into families
- Families follow well-known interaction paradigms





Some informal semantics of interaction paradigms



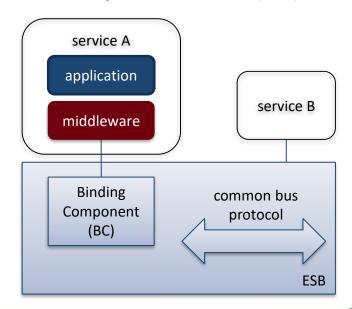
State-of-the-art in interaction paradigm interoperability

Theoretical approaches for individual paradigms

- Rely on concurrency theory, process algebras, architectural connectors
- No study of semantics *across interconnected paradigms*

Practical cross-paradigm approaches

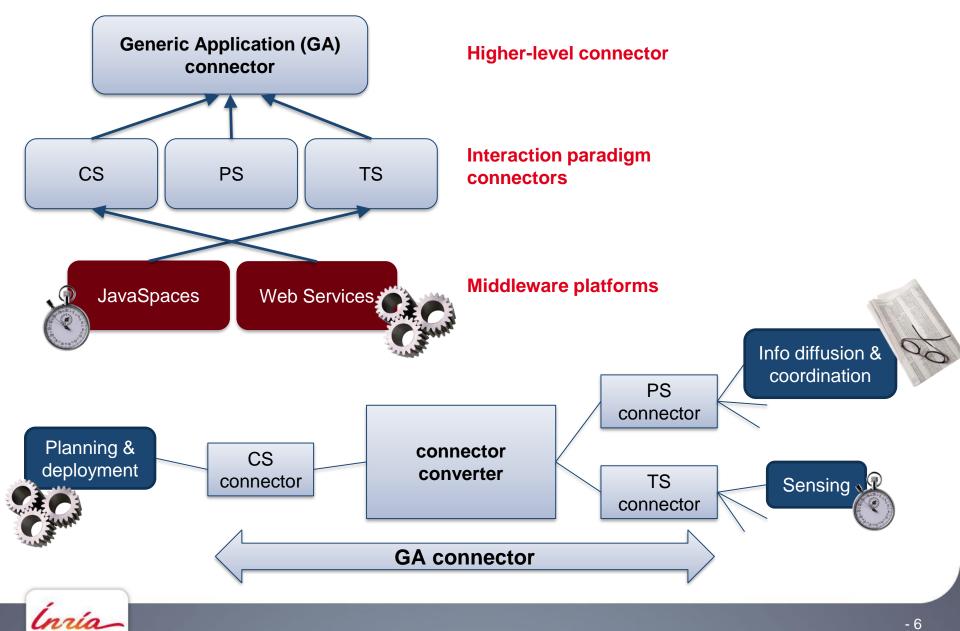
 Typically apply to specific middleware implementations







Solution to middleware interoperability

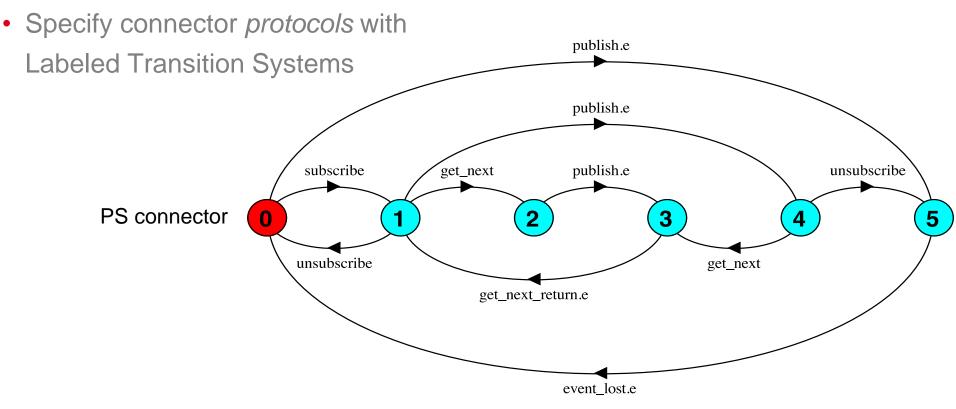


Mapping of space coupling – IDLs

	CS	PS	TS	GA
element	message	event	tuple	data
main scope	system ID	system ID	system ID	system ID
suk-scope	operation	topic	template	data qualifier
interaction semantics	{one-way, notification, request-response, solicit-response}	{publish, subscribe}	{write, take, read}	{post, get, post-get, get-post}
S&R element	<i>in:</i> {sensorLocation, lifeSign} <i>out:</i> {personnelLocation}	<i>out:</i> {personnelld, personnelLocation}	<i>out:</i> {sensorLocation, lifeSign}	in/out data



Time coupling and concurrency of CS, PS, TS connectors



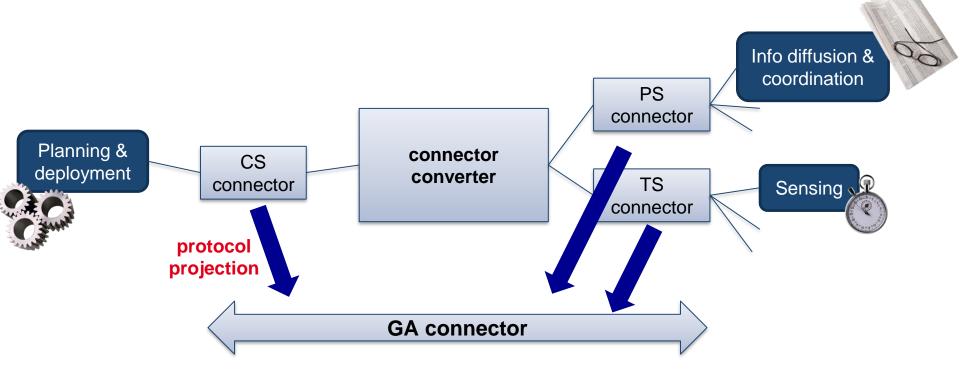
• Express and verify *semantics* in LTL temporal logic

fluent SUBSCRIBED = < {subscribe}, {unsubscribe} >

assert **EVENT_RECEIVED_IF_SUBSCRIBED_BEFORE_PUBLISH** = forall [e:EVENT] [] ((SUBSCRIBED && publish[e]) -> (!event_lost[e] U get_next_return[e]))



Mapping of time coupling and concurrency

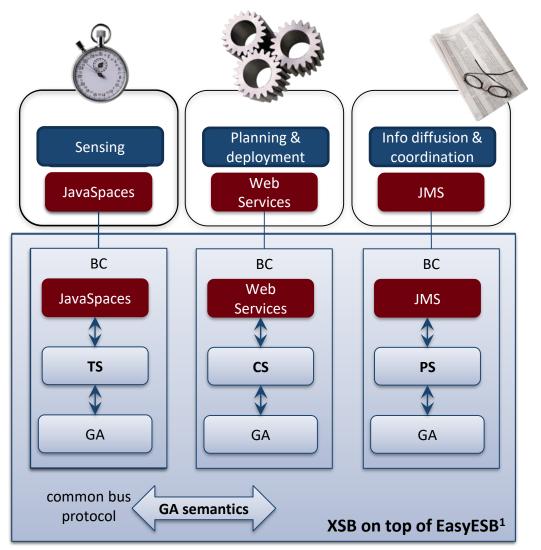


- Rely on the method of protocol conversion via projections¹
- Common semantics of CS, PS, TS apply to GA end-to-end

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Implementation of the GA connector: an eXtensible Service Bus (XSB)

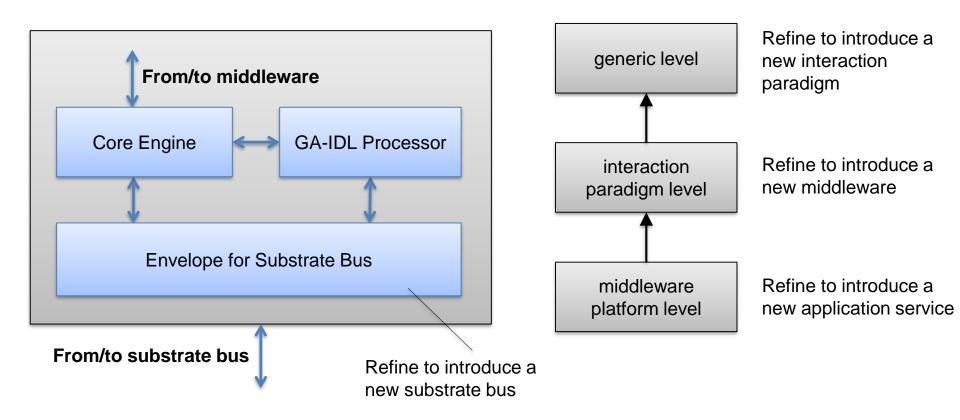


- XSB is a generic bus
- Build on top of a substrate bus
- Cross-paradigm integration in the S&R application



XSB architectural framework

BC extensible architecture





Evaluation – Effort for S&R application design

Development effort for composing the three available application services into the S&R application

	CS-, PS-, TS-IDL	Generated GA-IDL	Data mapping
	descriptions	descriptions	directives
	(XML lines)	(XML lines)	(XML lines)
All 3 application services	407	249	226

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Evaluation – Extensibility of the XSB framework

Development effort for the JMS binding component

	Lines of code	PS-, GA-IDL schemas (XML lines)	Configuration files (XML lines)
Total effort	18021	2836	248
Developer's effort	1162	191	12
Developer's effort ratio	6%	6%	4%

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Evaluation - Performance

- Latency overhead introduced by the bus for all interconnection combinations
- Comparison between XSB and substrate bus

Latency overhead (ms)	Comparison
258	-
261,5	+1,4%
283	+9,7%
276	+7,0%
298	+15,5%
	261,5 283 276

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Conclusion and future perspective

- Service-oriented applications in the Future Internet
 - Require cross-middleware interoperability
- Tackle this challenge via
 - Abstractions and mappings for interaction paradigms
 - Assessment of end-to-end semantics
 - Implementation into an extensible service bus
- Next step
 - Extend with support for continuous interactions data streaming protocols





Further information:

XSB: xsb.inria.fr CHOReOS: www.choreos.eu Inria ARLES: www.rocq.inria.fr/arles

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